

Lec 10

Recall

fun add (x, y) = x + y

[[] (fn (x, y) => x + y) / add]

Consider

fun plus x = fn y => x + y plus: int -> int -> int

[(fn x => fn y => x + y) / plus] *can be thought of as waiting for next arg*

val inc = plus 1

inc: int -> int

[[1/x] (fn y => x + y) / inc]

inc 4 => [1/x] (fn y => x + y) 4
=> [1/x, 4/y] (x + y)
=> 1 + 4
=> 5

Good compiler usually optimise this, so should be same as add

Or... we can do one line?

plus 1 4 => [1/x] (fn y => x + y) 4

↳ Another way to do function call!
plus 1 4 ≅ add 1 4

plus 1 4

↑
"currying" — name from Haskell Curry

Syntactic sugar for currying

fun plus x y = x + y

fun plus x = (fn y => x + y)

val plus = (fn x => (fn y => x + y))

fun sum x y z = x + y + z

sum : int -> int -> int -> int

Higher order function (HOF ...?)

Def: functions that takes in some function as arg or returns a HOF as output.

(* filter : ('a -> bool) -> 'a list -> 'a list

REC: p is total

ENS: filter p L => L' s.t. L' ⊆ L and ∀ l ∈ L, p(l), preserve order and multiplicity.

(sloppy)

*)

fun filter p [] = []

| filter p x::xs = if p x then x::filter p xs else filter p xs end

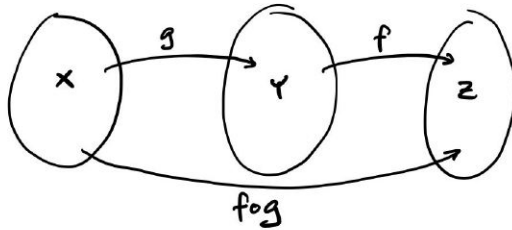
Note func app higher priority than infix ops

val keepevens = filter (fn n => n mod 2 = 0) : int list -> int list

keepevens [1,2,3,4,5] ↦ [2,4]

filter (fn n => n mod 2 = 0) [1,2,3,4,5] ↦ [2,4]

Composing function



already built-in.
 $\text{fun } (f \circ g)(x) = f(g(x))$

$o: ('b \rightarrow 'c) * ('a \rightarrow 'b) \rightarrow 'a \rightarrow 'c$

$\text{val increment} = \text{fn } x = x + 1$
 $\text{val double} = \text{fn } x = 2 * x$

$\text{increment} \circ \text{double} : \text{int} \rightarrow \text{int} \cong \text{fn } x \Rightarrow 2 * x + 1$
 $\text{double} \circ \text{increment} : \text{int} \rightarrow \text{int} \cong \text{fn } x \Rightarrow 2 * x + 2$

The most famous HOF — "map"

$(*) \text{ map} : ('a \rightarrow 'b) \rightarrow 'a \text{ list} \rightarrow 'b \text{ list}$

REQ true

ENS $\text{map } f [x_1, \dots, x_n] \cong [f x_1, \dots, f x_n]$

*)

$\text{fun map } f [] = []$

$| \text{map } f x::xs = (f x)::(\text{map } f xs)$

$\text{map double } [1, 2, 3] \rightarrow [2, 4, 6]$

$\text{val doubler} = \text{map double} : \text{int list} \rightarrow \text{int list}$

$\text{doubler } [1, 2, 3] \rightarrow [2, 4, 6]$

Fold : accumulate result over whole list

foldl, foldr : ('a * 'b → 'b) → 'b → 'a list → 'b

Def: foldr f z [x₁, ..., x_n] = f(x₁, ... f(x_{n-1}, f(x_n, z))...)

foldl f z [x₁, ..., x_n] = f(x_n, ... f(x₂, f(x₁, z))...)

foldl (op +) 0 [1, 2, 3, 4] ↪ 10

foldr (op +) 0 [1, 2, 3, 4] ↪ 10

foldl (op -) 0 [1, 2, 3, 4] ≅ 4 - (3 - (2 - (1 - 0))) ↪ 2

foldr (op -) 0 [1, 2, 3, 4] ≅ 1 - (2 - (3 - (4 - 0))) ↪ -2

fun foldl f z [] = []
| foldl f z x::xs = foldl f (f(x,z)) xs
↪ updated accumulator
↪ tail recursive

fun foldr f z [] = []
| foldr f z x::xs = f(x, foldr f z xs)
↪ which not tail recursive

ASIDE
foldr is
"catamorphism" for list.
see cat theory (?)
rather typical in data structs

foldr (op ::) [] ≅ id
foldl (op ::) [] ≅ rev
foldr (op ::) Y X ≅ X @ Y
foldl (op ::) acc L ≅ rev L @ acc