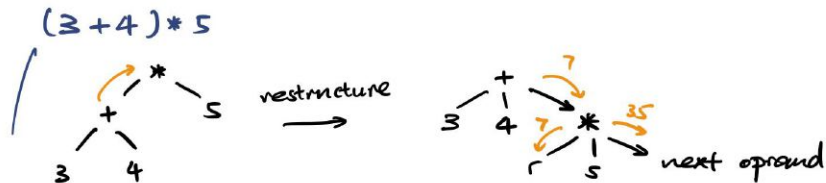


Lec 12

Continuation passing style (CPS) ↔ counterpart ↔ direct style

Consider:



$\text{fun add } x \ y \ k = k \ (x + y)$
 $\text{fun mult } x \ y \ k = k \ (x * y)$

$\text{int} \rightarrow \text{int} \rightarrow (\text{int} \rightarrow 'a) \rightarrow 'a$

$\text{add } 3 \ 4 \ (\text{fn } r \Rightarrow \text{mul } r \ 5 \ \underline{\text{Int.toString}})$

* answer type string
 * initial continuation

continuation
 output type from continuation
 whatever to do next

$\Rightarrow (\text{fn } r \Rightarrow \text{mul } r \ 5 \ \text{Int.toString}) (3+4)$
 $\Rightarrow (\text{fn } r \Rightarrow \text{mul } r \ 5 \ \text{Int.toString}) \ 7$
 $\Rightarrow [7/r] (\text{fn } r \Rightarrow \text{mul } r \ 5 \ \text{Int.toString})$
 $\Rightarrow [7/r] (\text{mul } 7 \ 5 \ \text{Int.toString})$
 $\Rightarrow [7/r] (\text{Int.toString } (7*5))$
 $\Rightarrow \text{Int.toString } 35$
 $\Rightarrow "35"$

Ex in different styles

```
(* sum : int list → int *)  
fun sum [] = 0  
  | sum x::xs = x + sum xs
```

```
fun sum L = foldl (op +) 0 L
```

```
val sum = fold (op +) 0
```

```
(* tsum : int list * int → int  
   ENS tsum (L, acc) ≡ sum L + acc  
  *)  
fun tsum ([], acc) = acc  
  | tsum (x::xs, acc) = tsum (xs, x + acc)  
fun sum L = tsum (L, 0)
```

Familiar ways to
write sum

↙ CPS

```
(* csum : int list → (int → 'a) → 'a
```

```
  REQ true
```

```
  ENS csum L k ≡ k (sum L)
```

```
  *)
```

```
fun csum [] k = k 0
```

```
  | csum x::xs k = csum xs (fn r ⇒ k (x + r))
```

```
fun sum L = csum L (fn x ⇒ x)
```

Hmm everything's tail call, but
not a way to improve performance.
Compiler turns normal sum into
CPS!

But notice we're building longer
and longer k?

Those inferior languages (...?)

C doesn't pretend to be a programming language
... but python is not a real programming language
but it pretends to be one...

← random digression / joke

Exercise : trace csum [2,3] Int.toString

Back to trees

`datatype 'a tree = Empty | Node of 'a tree * 'a * 'a tree`

`(* 'a tree → 'a list → 'a list *)`

`fun inorder Empty acc = acc
 | inorder (Node (L, x, R)) acc = inorder L (x :: inorder R acc)`

`(* same int list → int list → bool *)`

`fun same [] [] = true
 | same x::xs y::ys = x = y andalso same xs ys
 | same _ _ = false` ↑ able to stop early

`(* treematch : int tree → int list → bool *)`

`fun treematch T L = same (inorder T []) L`

⚠ bad performance if inorder before any comparison

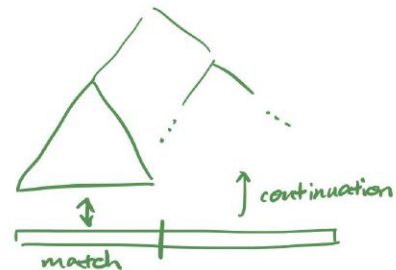
CPS version — capable of early stopping

`(* prefix : int tree → int list → (int list → bool) → bool
 REQ k is total
 ENS prefix T L k ⇒ { true if L ≡ L1 @ L2 s.t. L1 is inord of T
 and k L2 ⇒ true
 false otherwise
 *)`

`fun treematch T L = prefix T L (fn [] ⇒ true | _ ⇒ false)`
└ built in as List.null ┘ └ initial continuation ┘

`fun prefix (Empty : int tree) L k = k L`

`| prefix (Node (l, x, r)) L k = prefix l L (fn [] ⇒ false |
↑ tail recursive! └ this is not call, just passed in as value ┘
 y::ys ⇒ x = y andalso prefix r ys k)`



Ex. Polymorphise (...?)

Two continuations !

(^{cont.}* search (^{predicate}'a → bool) → 'a tree → (^{cont. success-callback}'a → 'b) → (^{cont. failure-callback}unit 'b) → 'b
REQ p total
⋮

*)

```
fun search p Empty sc fc = fc ()  
  | search p (Node(L, x, R)) sc fc =  
    if p x then sc x  
    else  
      search p L sc (fn () ⇒ search p R sc fc)
```

```
fun findeven = search (fn n ⇒ n mod 2 = 0) T  
  (fn x ⇒ SOME x)  
  (fn () ⇒ NONE)
```