

# Lec 16

Today: most important concept in CS — Abstraction!

Story: working on Excel — modular, but fully connected. Hard to abstract.  
everything depend on everything, so little protection from other programmers  
(you'll be working with idiots)

# Enforcing abstraction in ML

Concept		ML
interface	—	signature
implementation	—	structure / module

↓ expression

Int. compare

↑ module

signature ↓

Int : INTEGER

## Signature

Signature QUEUE =

Sig

type 'a q ← abstract! Client don't know what it equals to

val empty: 'a q

val enq: 'a q \* 'a → 'a q

val null: 'a q → bool

val deq: 'a q → 'a \* 'a q

exception Empty

end

## Structure with list as queue

### 1. Abstraction function

The list contains all current elements of queue in arrival order

### 2. Representation invariant

None in this case

### 3. Code

Structure listQueue :> QUEUE =

Struct

type 'a q = 'a list

val empty = []

fun enq (q, x) = q @ [x]

fun null [] = true

| null \_ = false

exception Empty

fun deq [] = raise Empty

| deq x::q = (x, q)

end

opaque ascription, but can make transparent by replace ":>" with ":".  
↑

should only use for debugging

abstract type

need to match signature

can write things other than those in signature

Structure Q = ListQueue

## Client

val q2 = Q.enq (Q.enq (Q.empty, 1), 2)

q2 : int Q.q ← Client not allowed to treat it as int list

= - : int Q.q ← SMLNJ prints this

val (a, b) = Q.deq q2

val (c, \_) = Q.deq q2

val (d, \_) = Q.deq b

\* ML by default doesn't print what you're supposed to see.

## # Improving queue with double stack

### 1. Abstraction

(front, back) — front @ rev back contains elems of queue in arrival order

### 3. Code

Structure TwoListQueue :> QUEUE =

Struct

type 'a q = 'a list \* 'a list

val empty = ([], [])

fun enq (cf, b), x = (f, x::b)

← O(1)

fun null ([], []) = true

! null \_ = false

exception Empty

fun deq ([], []) = raise Empty

! deq ([], b) = deq (rev b, [])

← Better amortised O(1)

! deq (x::f, b) = (x, cf, b)

end

Structure Q = TwoListQueue

## # Dictionary

Signature DICT =

Sig

type key = string

concrete type

type 'a entry = key \* 'a

type 'a dict

abstract

val empty : 'a dict

val lookup : 'a dict → key → 'a option

val insert : 'a dict → 'a entry → 'a dict

end

Bad impl - list as dict [omit]

Better impl - tree

1. Abstraction function ← refers to mapping btwn abstract repr and theoretical data  
(key, value) pairs in tree correspond to entries in dict
2. Repr invar  
tree is sorted
3. Impl

Structure BST :> DICT =

Struct

type key = string

type 'a entry = key \* 'a

datatype 'a tree = Empty | Node of 'a tree \* 'a entry \* 'a tree

type 'a dict = 'a tree

val empty = Empty

fun insert (Empty, e) = Node (Empty, e, Empty)

| insert (Node (L, e' as (k', -), R), e as (k, -)) =

(case String.compare (k, k') of

EQUAL ⇒ Node (L, e, R)

LESS ⇒ Node (insert (L, e), e', R)

GREATER ⇒ Node (L, e', insert (R, e))

)

⋮

end

New field not in signature.  
But if transparent, it becomes  
broken as invisible fields in sig  
can get printed.

could define  
as datatype  
in struct

⚠ Don't pattern match structure