

## Lec 20

### Lazy

\* Dealing with infinite data structure.

Consider  $f$  vs  $(\lambda x \Rightarrow f x)$ . Are they the same?

→  $f$  could be expression, not value, but  $(\lambda x \Rightarrow f x)$  can only be value

$\text{fun loop } x = \text{loop } x$

$\text{loop } () \leftarrow \text{loops}$        $(\lambda x \Rightarrow (\text{loop } ()) x) \leftarrow \text{just waits there}$

### # Suspension

Def A suspension of type  $t$  is a func of type  $\text{unit} \rightarrow t$

Let  $e : t$      $\frac{(\lambda x \Rightarrow e) : \text{unit} \rightarrow t}{\text{any expression}}$  is a " $t$  suspension"

"lazy" representation of  $e$

Calling a suspension on  $()$  is called "forcing" the suspension

Def A stream is some infinite data structure.

### # Sig for stream

Signature STREAM =

sig

type 'a stream (\* abstract \*)

datatype 'a front = Empty | Cons of 'a \* 'a stream

val expose : 'a stream  $\rightarrow$  'a front

val delay : ( $\text{unit} \rightarrow 'a \text{ front}$ )  $\rightarrow$  'a stream

```

val empty : 'a stream
val null : 'a stream → bool
val take : 'a stream * int → 'a list
val map : ('a → 'b) → 'a stream → 'b stream
val filter : ('a → bool) → 'a stream → 'a stream
:
end

```

Structure Stream :> STREAM =

```

struct
  datatype 'a stream = Stream of (unit → 'a front)
  and 'a front = Empty | Cons of 'a * 'a stream
  fun expose (Stream d) = d()
  fun delay d = Stream d
  val empty = delay (fn () => Empty)
  fun null s = (case expose s of Empty => true | _ => false)
    could loop!
:
end

```

Structure S = Stream

Example usage

```

fun ones' () = S.Cons (1, S.delay (fn () => one' ()))      : ones' : unit → int S.front
val ones = S.delay (fn () => ones' ())                      : int S.stream

```

```

fun natsFrom' x = S.Cons (x, S.delay (fn () => natsFrom' (x+1)))
fun natsFrom x = S.delay (fn () => natsFrom' x)
val nats = natsFrom 0

```

## # More lemmas

Def Streams  $X, Y$  are eq if

$$S.\text{take}(X, n) \cong S.\text{take}(Y, n) \text{ for all } n$$

Def Stream  $X$  is productive if  $S.\text{expose } X$  returns  
1. Empty or  
2.  $\text{Cons}(x, X')$  with productive  $X'$

Def Stream  $X$  is infinite if it's productive and we never get Empty

## # Map

(\* map: ('a → 'b) → ('a Stream) → ('b Stream) \*)  
fun map f s = delay (fn () => map' f (expose s))  
and map' f Empty = Empty  
| map' f (Cons(x, s')) = Cons(f x, map f s')

map': ('a → 'b) → 'a front → 'b front

(\* filter : ('a → bool) → 'a Stream → 'a Stream \*)  
fun filter p s = delay (fn () => filter' p (expose s))  
and filter' p Empty = Empty  
| filter' p (Cons(x, s')) =  
 if p x then  
 Cons(x, filter p s')  
 else  
 filter' p (expose s')

## Example usage

```
val evens = S.map (fn x => 2*x) nats
val [0,2,4] = S.take (evens, 3)
val ns = S.filter (fn x => x < 0) nats   ← almost instantaneous
val _ = S.expose ns   ← loops forever.
                                                Notice ns is empty stream
                                                but the code cannot tell.
```

# Code to compute all prime



```
fun sieve s = S.delay (fn () => sieve' (S.expose))
and sieve' S.Empty = S.Empty
  | sieve' (S.Cons (p, s)) =
      S.Cons (p, sieve' (S.filter (fn x => x mod p ≠ 0) s))
```

val primes = sieve (natsFrom 2)