Lec 26 Memoisation with parallelism Things that can show up in 15-418, 15-312, 15-410 # Sequential impl

from memoise f = let val conche = ref Table.empty from g a = case Table.find ! cache a of SOME $r \Rightarrow r$ | NONE \Rightarrow let val r = f g a] \leftarrow multi-threads can call f if they have val r = f g a f same a and a not in table val f = cache := Table.insert cache f cache f at same time

in end

Prob we often do > 2 recursive calls and want to do them in parallel, but this impl not safe for parallelism

Idea suspend execution, make sure to not race compute

- at \triangle , insert busy marker, at \square , update actual result so lookup result can be busy, some, none.

- at , handle busy case by

- busy wait
 sleep wait (OS could schedule some other work?)
- suspend job (SML bridt-in, but not Python 1C+, etc.)
 give continuotion / houndle +0 another +hread

 SML: callect to suspend
 throw to wake up

 The put self in some queue
 try wake things up at a

Impl initialise empty guene Q states for table entry state = wait of Q I full of B

fun $g = case Table find ! cache a of NONE <math>\Rightarrow$ insert (a, wait (empty queue)) to table; r = fga;

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get queue Q;
insert (a, full r);
wake up everything in Q ] < OCI) span with some impl

| SOME (want Q) => cuspend self in Q that supports punalled map

| SOME (full r) => r

Also, make sure table and queue are linearisable
concurrent data struct needed.

# Concurrent table

insert: ctable > (x x p) > p option

T (a, b)

if a not in T, add (a, b) and return NONE

if (a, b') in T, return SOMEb' <- can retry with update
in this case

update: ctable > (x, p) > ()

let

cache = Table.empty
fun g a =

let

Q = Queue.empty
in

case (Ctable.insert cache (a, wait Q)) of
```

SOME (full r) ⇒ r I SOME (wait (2) ⇒

| NONE ⇒ ! (as before)

suspend self onto Q; when wake up, find result and return